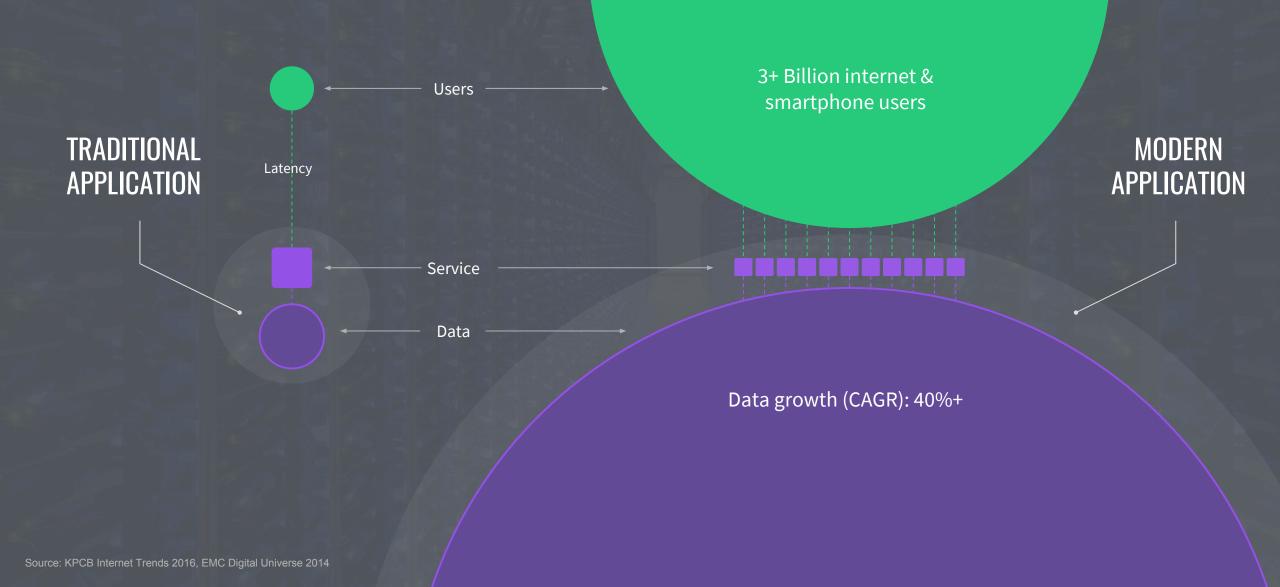
June 2017

DC/OS AND FAST DATA (THE SMACK STACK)



Benjamin Hindman - @benh Elizabeth K. Joseph - @pleia2

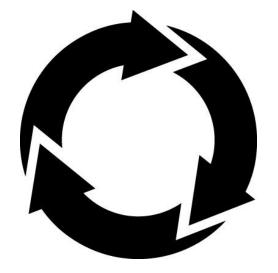
ARCHITECTURAL SHIFT



TODAY'S REINFORCING TRENDS

CONTAINERIZATION

MICROSERVICES



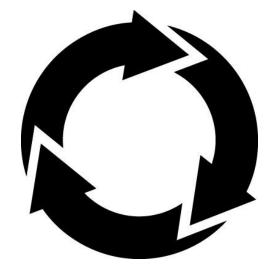
CONTAINER ORCHESTRATION

BIG DATA & ANALYTICS

TODAY'S REINFORCING TRENDS

CONTAINERIZATION

MICROSERVICES



CONTAINER ORCHESTRATION

FAST BIG DATA & ANALYTICS

FROM BIG DATA TO FAST DATA

Days	Hours Minu		s Seconds		Microseconds			
Ba	atch		Micro-Batch		Event Pr	ocessing		
Reports wha	t has happened usin	g descriptive a	nalytics	Solves p	Solves problems using predictive and prescriptive analytics			
Billing, Chargeb	ack Produc	t recommend	dations Real-time Pricing	g and Routing	Real-time Advertising	Predictive User Interface		
	. a	ers Who Bought This Item Also Boug						





HDMI Cable - 6.5 Feet (2 Curved Panel UHD HD TV Wall Mount Bracket Meters) Supports Ethernet for 32-65* TVs (Many from Articulating Arm Swivel & 3D, 4K and Audio Return 20-75") Specifically fits up







5

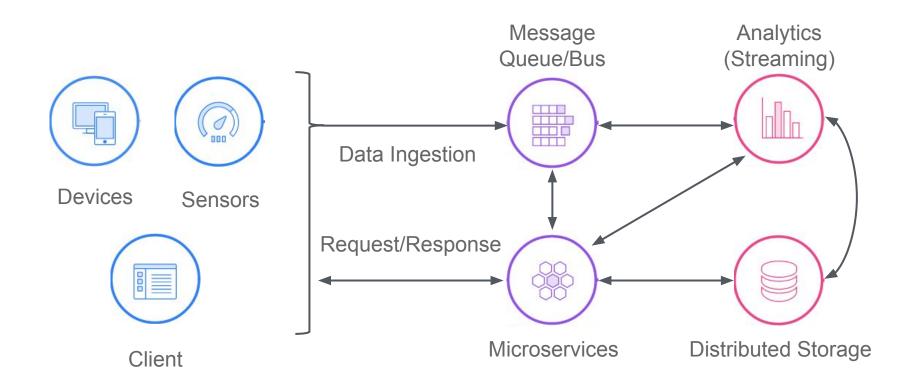
ON THE EDGE, AND STILL REALLY BIG!

A380-1000: 10,000 sensors in each wing; produces more than 7Tb of IoT data per day

[1] https://goo.gl/2S4q5N

6

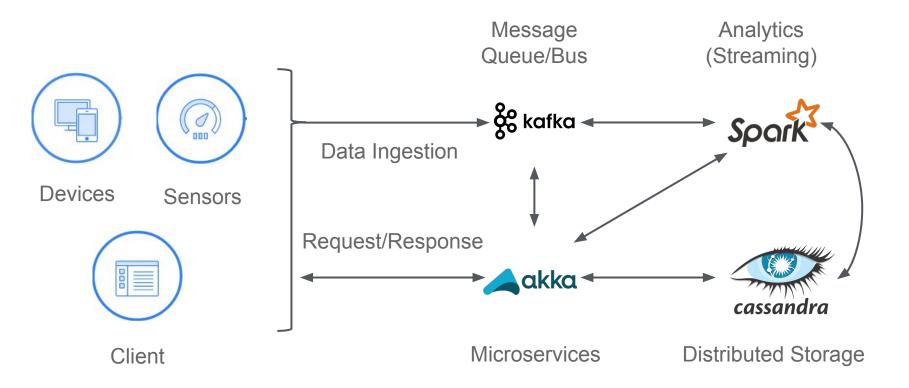
MODERN APPLICATION -> FAST DATA BUILT-IN



Use Cases:

- Anomaly detection
- Personalization
- IoT Applications
- Predictive Analytics
- Machine Learning

A GOOD STACK ...



Use Cases:

- Anomaly detection
- Personalization
- IoT Applications
- Predictive Analytics
- Machine Learning

MESSAGE QUEUES



Message Brokers

- Apache Kafka
- ØMQ, RabbitMQ, Disque

Log-based Queues

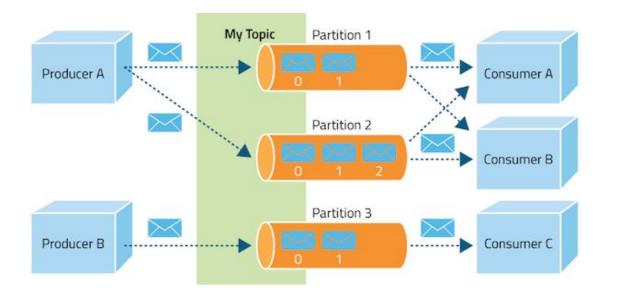
• fluentd, Logstash, Flume

see also <u>queues.io</u>



9

APACHE KAFKA



Typical Use: A reliable buffer for stream processing

Why Kafka?

- High-throughput, distributed, persistent publish-subscribe messaging system
- Created by LinkedIn; used in production by 100+ web-scale companies [1]

[1] https://cwiki.apache.org/confluence/display/KAFKA/Powered+By

DELIVERY GUARANTEES

- At most once—Messages may be lost but are never redelivered
- At least once—Messages are never lost but may be redelivered (Kafka)
- **Exactly once**—Messages are delivered once and only once (this is what everyone actually wants, but no one can deliver!)

Murphy's Law of Distributed Systems:

Anything that can go wrong, will go wrong ... partially!

STREAMING ANALYTICS

Microbatching

• Apache Spark (Streaming)

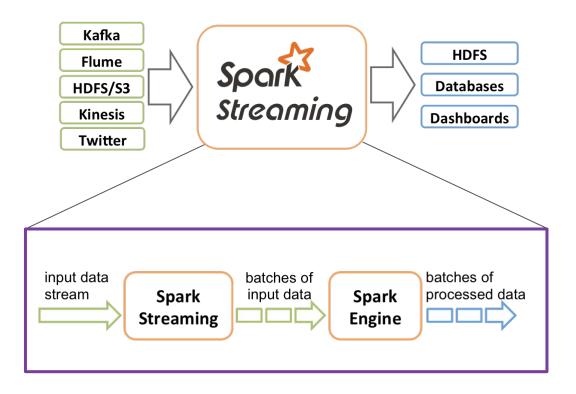
Native Streaming

- Apache Flink
- Apache Storm/Heron
- Apache Apex
- Apache Samza





APACHE SPARK (STREAMING)



Typical Use: distributed, large-scale data processing; micro-batching

Why Spark Streaming?

- Micro-batching creates very low latency, which can be faster
- Well defined role means it fits in well with other pieces of the pipeline

DISTRIBUTED STORAGE

NoSQL

- ArangoDB
- mongoDB
- Apache Cassandra
- Apache HBase

SQL

MemSQL

Filesystems

- Quobyte
- HDFS

Time-Series Datastores

- InfluxDB
- OpenTSDB
- KairosDB
- Prometheus

see also iot-a.info







APACHE CASSANDRA



Typical Use: No-dependency, time series database

Why Cassandra?

- A top level Apache project born at Facebook and built on Amazon's Dynamo and Google's BigTable
- Offers continuous availability, linear scale performance, operational simplicity and easy data distribution

how do we operate these distributed systems?

most organizations have many stateless independent (micro)services, the *distributed systems* I'm talking about here are ...

latency state consensus synchronous $\mathbf{\sigma}$ fault cation topology ISTIC etermin D as theorem rmance scalability bandwidth power insecure stateful multiple

how do we scale the operations of distributed systems?

SMACK STACK





Apache Spark: distributed, large-scale data processing

Apache Mesos: cluster resource manager



Akka: toolkit for message driven applications



Apache Cassandra: distributed, highly-available database



Apache Kafka: distributed, highly-available messaging system

distributed systems are hard to operate

DATA & ANALYTICS DAY 2 OPS CHALLENGES

- Bare metal storage (or someone else's problem)
- Drive down job latency and drive up utilization
- Run multiple versions simultaneously
- Upgrade complicated data systems

1: download 2: deploy 3: monitor 4: maintain 5: upgrade → goto 1

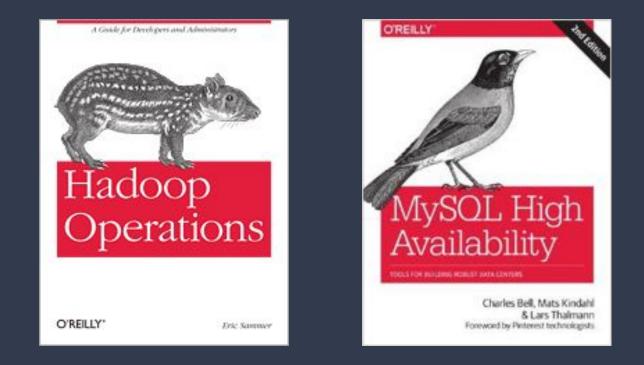
fault tolerance

- + high availability
- + latency

╋

- + bandwidth
- + CPU/mem resources

= configuration



#!/bin/bash

pip install "\$1" & easy_install "\$1" & brew install "\$1" & npm install "\$1" & docker run "\$1" & http://www.stall "\$1" & apt-get install "\$1" & sudo apt-get install "\$1" & steamcmd +app_update "\$1" validate & git clone https://github.com/"\$1"/"\$1" & cd "\$1";./configure;make;make install & curl "\$1" | bash &



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Puppet A A

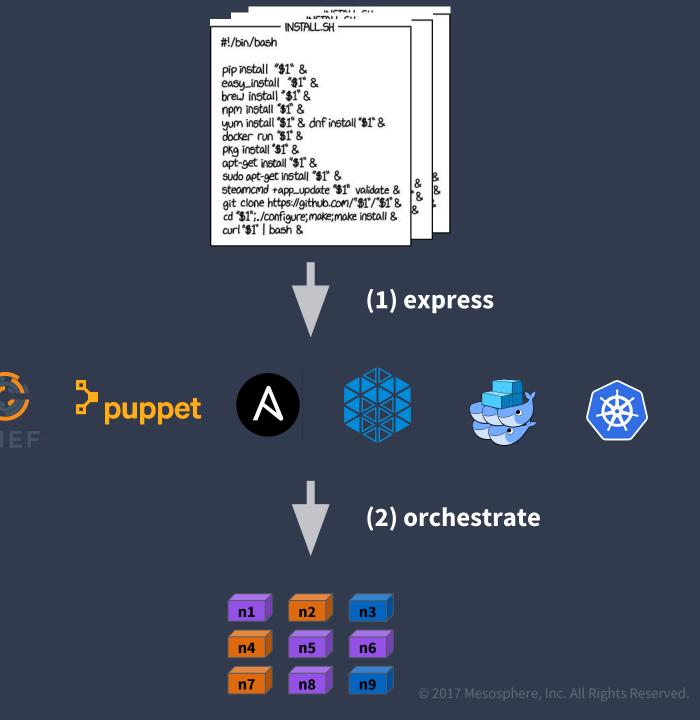
(1) express

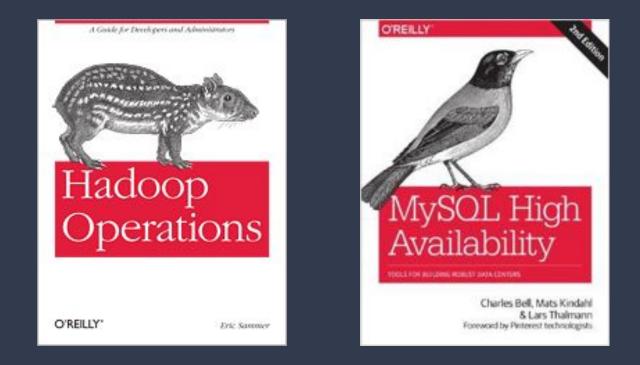


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pip install "\$1" & easy_install "\$1" & brew install "\$1" & yum install "\$1" & docker run "\$1" & http://www.stall pkg install "\$1" & apt-get install "\$1" & sudo apt-get install "\$1" & steamcmd +app_update "\$1" validate & git clone https://github.com/"\$1"/"\$1" & cd "\$1"; /configure; make; make install & curl "\$1" | bash &

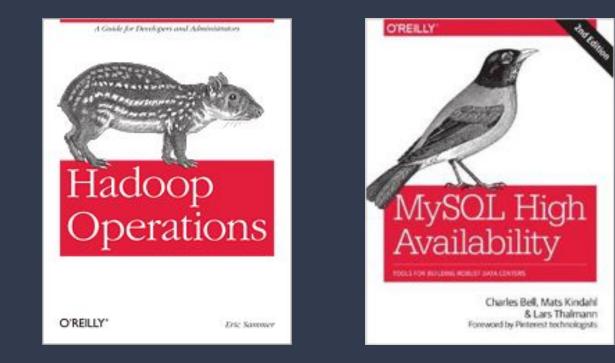
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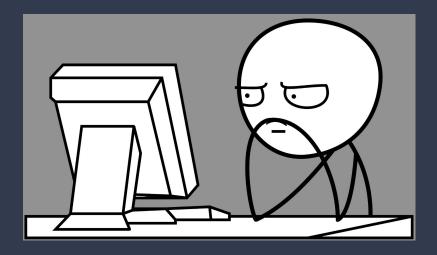


Nagios General Home Documentation Monitoring Tactical Overview	Current Network Status Last Updated: Sun Jul 15 14:03:12 CDT 2001 Updated every 75 seconds Nagios TM - <u>unwww.nagiosong</u> Logged in as guest - Monitoring process is running - Notifications cannot be sent out! - Service checks are being executed			<u>Up</u> 28	3 4		Service Status TotalsOkWarningUnknownCriticalPending102201418All ProblemsAll Types18137	0
 Status Detail Status Overview Status Summary Status Grid Status Map 3-D Status Map Service Problems Network Outages 	Display Fil Host Status Ty Host Propertie	ions For All Hosts Iters: /pes: All /s: Any s Types: All Problems		Se	rvice Details F	or All Hosts		B
Trends Availability	Host 个	Service 1	Status 个	Last Check 个	Duration 🚹	Attempt 🕆 🖡	Service Information	
Alert History	bogus-router	PING	CRITICAL	07-15-2001 13:59:39	4d 3h 43m 17s	1/3	CRITICAL - Plugin timed out after 10 seconds	
Notifications Log File	bogus1	Something	CRITICAL	07-15-2001 14:00:38	4d 3h 58m 49s	1/3	(Service Check Timed Out)	
© Comments		PING	CRITICAL	07-15-2001 14:02:36	4d 3h 58m 49s	1/3	CRITICAL - Plugin timed out after 10 seconds	
Downtime	bogus2	PING	CRITICAL	07-15-2001 13:59:09	4d 3h 44m 27s	1/3	CRITICAL - Plugin timed out after 10 seconds	
Process Info		Something	CRITICAL	07-15-2001 13:59:39	4d 3h 42m 26s	1/3	(Service Check Timed Out)	
Performance Info	bogus3	PING	CRITICAL	07-15-2001 14:00:38	4d 3h 42m 7s	1/3	CRITICAL - Plugin timed out after 10 seconds	
Configuration		Something	CRITICAL	07-15-2001 13:57:36	4d 3h 30m 35s	1/3	(Service Check Timed Out)	
View Config	bogus4	PING	CRITICAL	07-15-2001 13:59:09	4d 3h 43m 35s	1/3	CRITICAL - Plugin timed out after 10 seconds	
	and a second sec	Something	CRITICAL	07-15-2001 13:59:39	4d 3h 42m 26s	1/3	(Service Check Timed Out)	
	bogus5	PING	CRITICAL	07-15-2001 14:00:43	4d 3h 41m 7s	1/3	CRITICAL - Plugin timed out after 10 seconds	
	00000	Something	CRITICAL	07-15-2001 13:57:36	4d 3h 41m 7s	1/3	(Service Check Timed Out)	

first, debug ...



first, debug ...

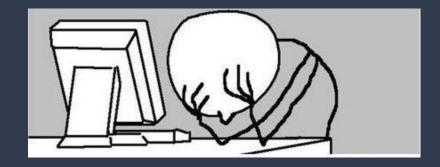


second, fix (scale, patch, etc)

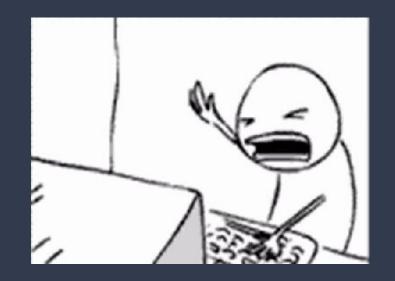
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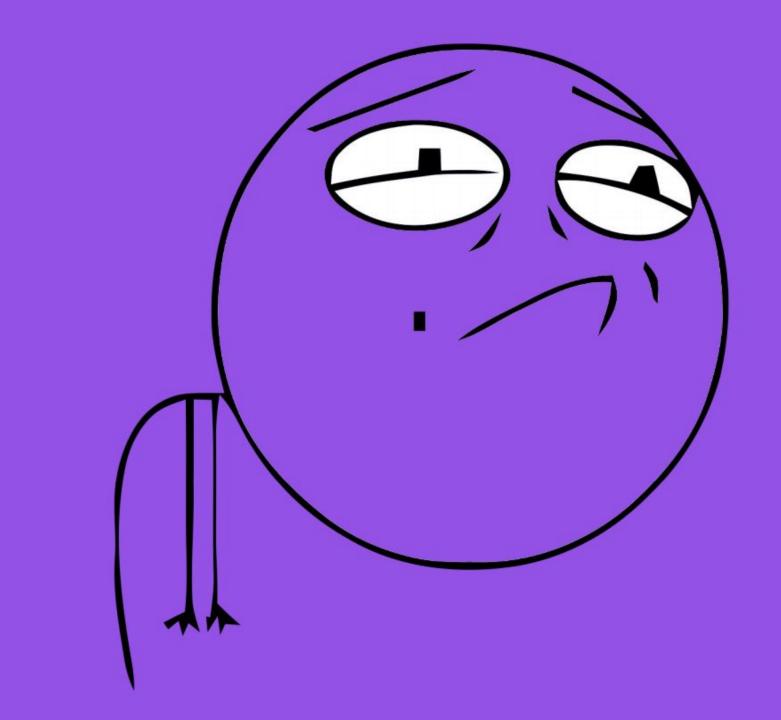
then, debug again ...



finally, write code so it never happens again ...



1: download
 2: deploy
 3: monitor
 4: maintain
 5: upgrade → goto 1



thesis: distributed systems should (be able to) operate themselves; deploy, monitor, upgrade ...

why: (1) operators have *inadequate* knowledge of distributed system needs/semantics to make optimal decisions

why: (1) operators have inadequate knowledge of distributed system needs/semantics to make optimal decisions (even after reading the book)

why: (2) execution needs/semantics can't easily or efficiently be expressed to underlying system, and vice versa





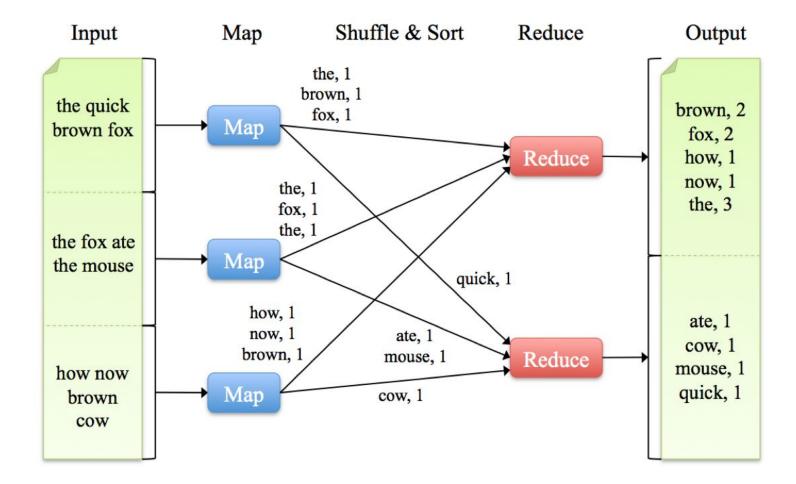
#!/bin/bash

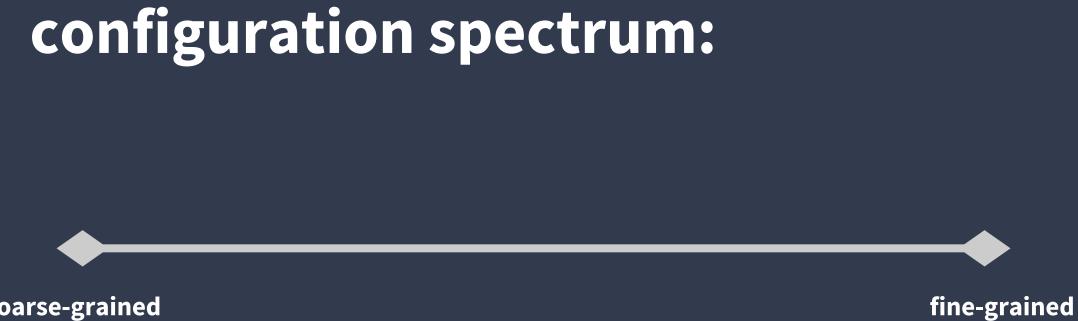
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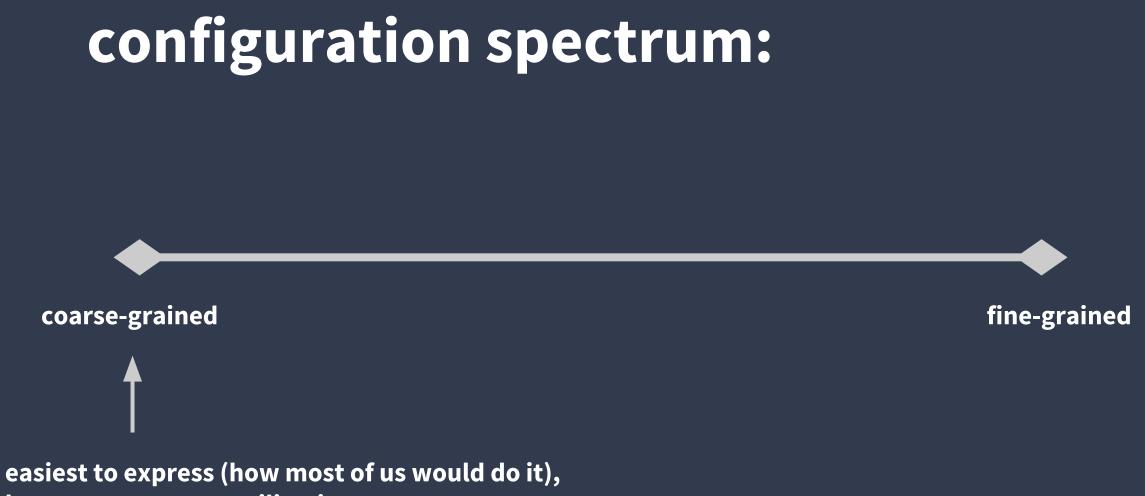
(2) orchestrate







coarse-grained

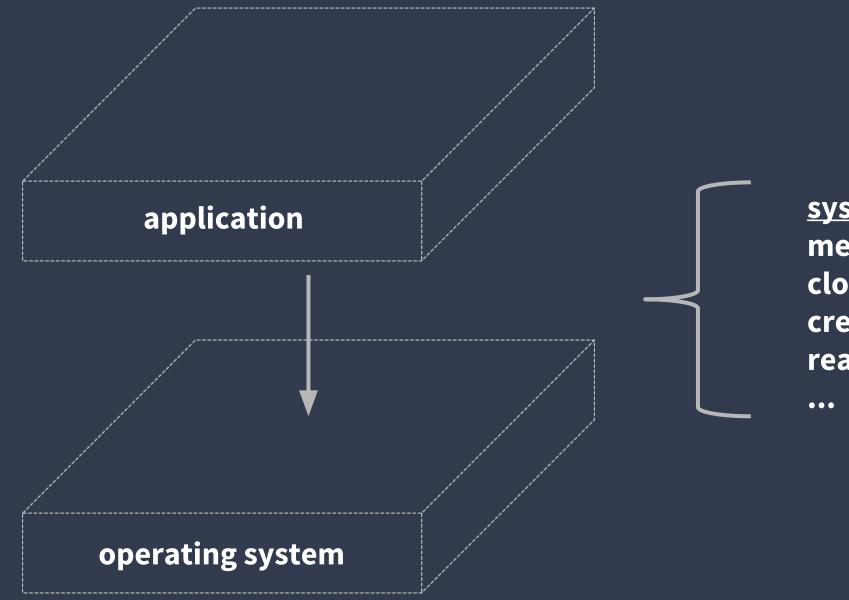


but worst resource utilization



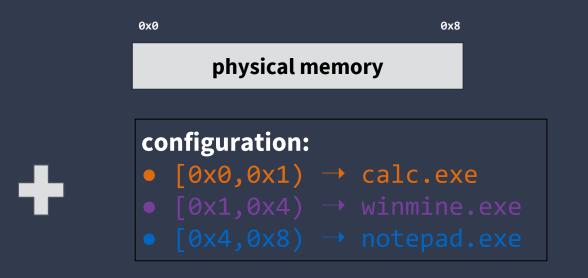
why can't Hadoop decide this for me?

applications "operate" themselves on Linux; when an application needs to "scale up" it asks the operating system to allocate more memory or create another thread ...



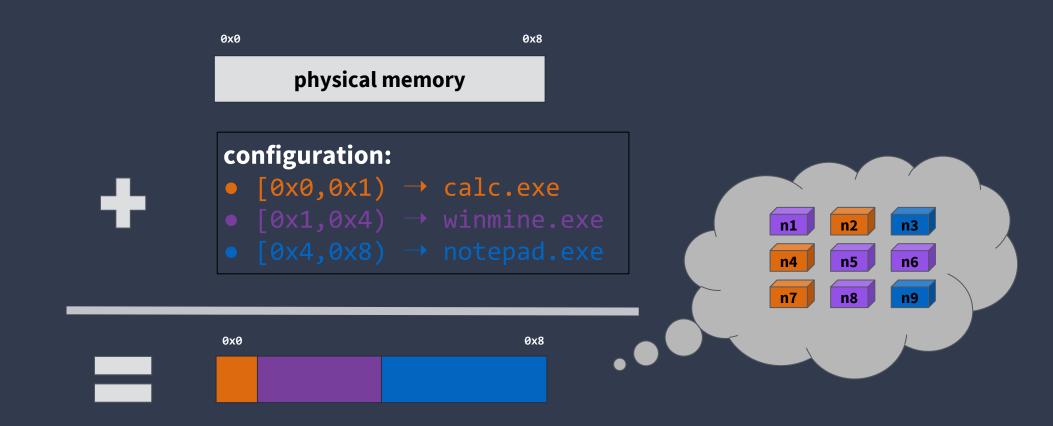
syscall interface: memory allocate clone/fork create file read, write

once upon a time ... before virtual memory

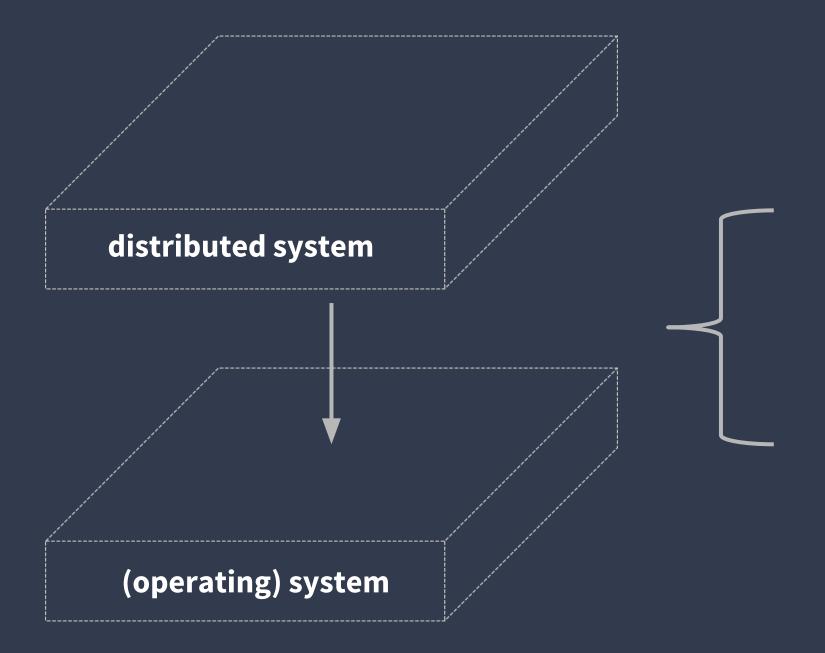




once upon a time ... before virtual memory



how: distributed systems need interface to communicate with underlying system, and vice versa

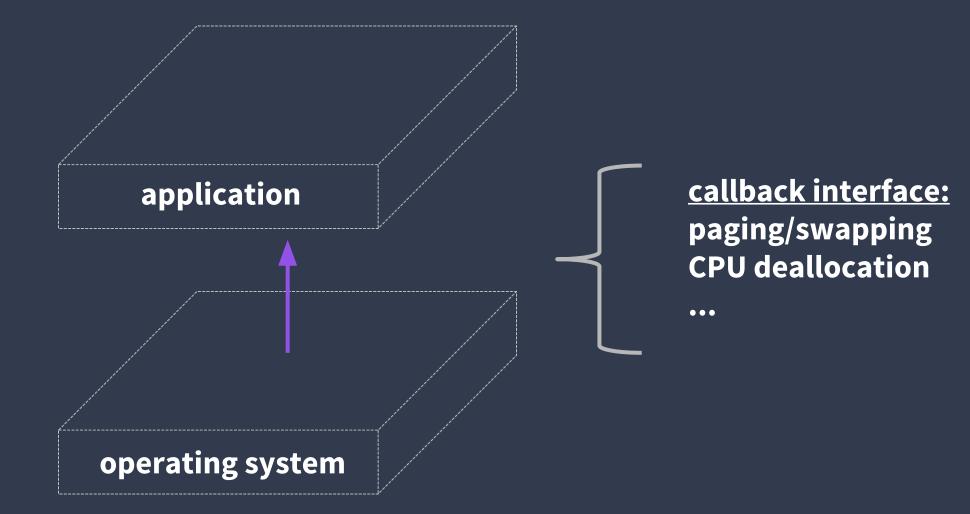


interface: resource allocation launch container/VM create storage attach/detach storage

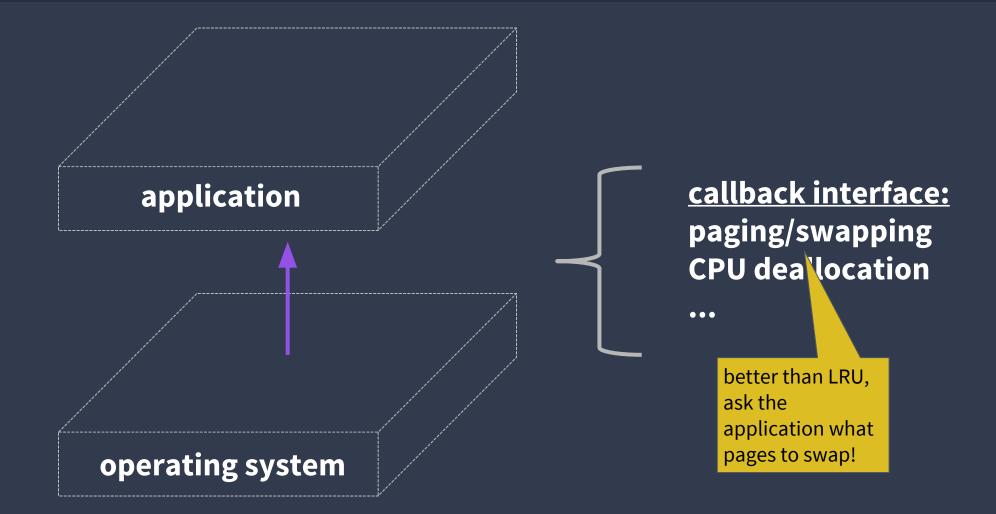
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vice versa: operating system should be able to callback into application

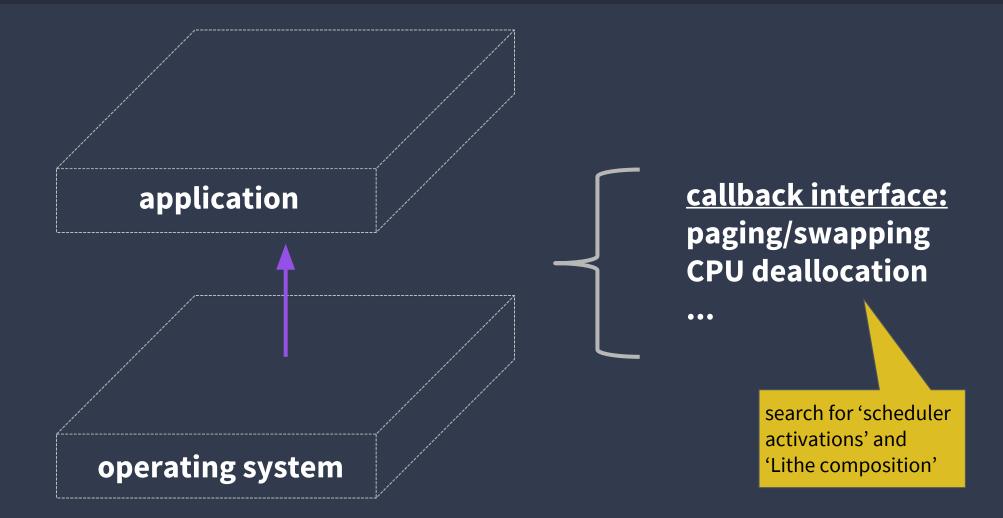
learning from history ... bidirectional interface

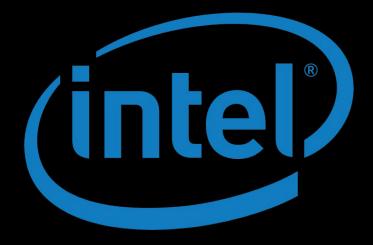


learning from history ... bidirectional interface



learning from history ... bidirectional interface







Enable MKL threading - use when you are sure that there are enough resources (physical cores) for MKL threading in addition to your own threads. Choose N carefully.

Example 1:

application has 2 threads, each thread calls MKL and the system has 8 cores: it's reasonable to set MKL_NUM_THREADS=4.

Example 2:

MKL function is called from a critical section of a parallel region - set MKL_NUM_THREADS=N, where N is the number of physical cores in the system (or use mkl_set_num_thread(N) routine).

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Software Products

Intel® Math Kernel Library (Intel® MKL) Using Intel® MKL with Threaded Applications

Page Contents:

- Memory Allocation MKL: Memory appears to be allocated and not released when calling some Intel MKL routines (e.g. sgetrf).
- Using Threading with BLAS and LAPACK
- Setting the Number of Threads for OpenMP (OMP)
- Changing the Number of Processors for Threading During Runtime
- Can I use Intel MKL if I thread my application?

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One of the advantages of using the IntelMKL is that it is multithreaded using OpenMP*. OpenMP* requires buffers to perform some operations and allocates even for single-processor systems and single-thread applications with memory occurs once the first time the OpenMP software is encountered, the program, allocation persists until the application terminates. In addition, the Windows* ope will allocate a stack equal to the main stack for every additional thread created, of memory that is automatically allocated will depend on the main stack, the Ope allocations and the number of threads used.

Using Threading with BLAS and LAPACK

Intel MKL is threaded in a number of place LAPACK ("GETRF, "POTRF, "GB Level 3 BLAS, DFTs, and FFTs. Intel M uses OpenMP" threading software. T situations in which conflicts can exist that make the use of threads in Intel MKL p We list them here with recommendations for dealing with these. First, a brief dis the problem exists is appropriat

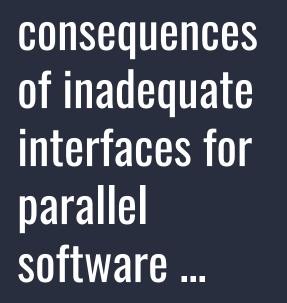
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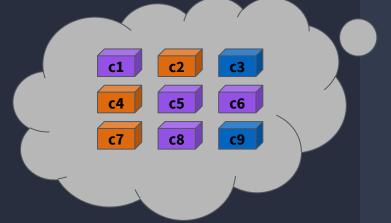
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 User threads the program using OpenMP directives and/or pragmas and compiles the program using a compiler other than a compiler from Intel. This is more problematic because setting OMP_NUM_THREADS in the environment affects both the compiler's threading library and the threading

http://www.intel.com/support/performancetools/libraries/mkl/sb/CS-017177.htm

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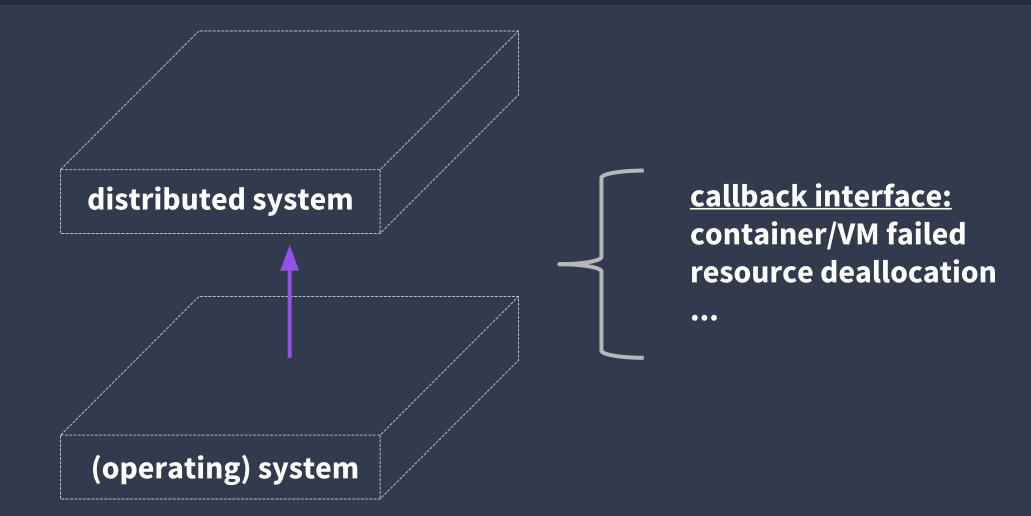
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operating system has inadequate knowledge of applications execution needs/semantics to make optimal decisions

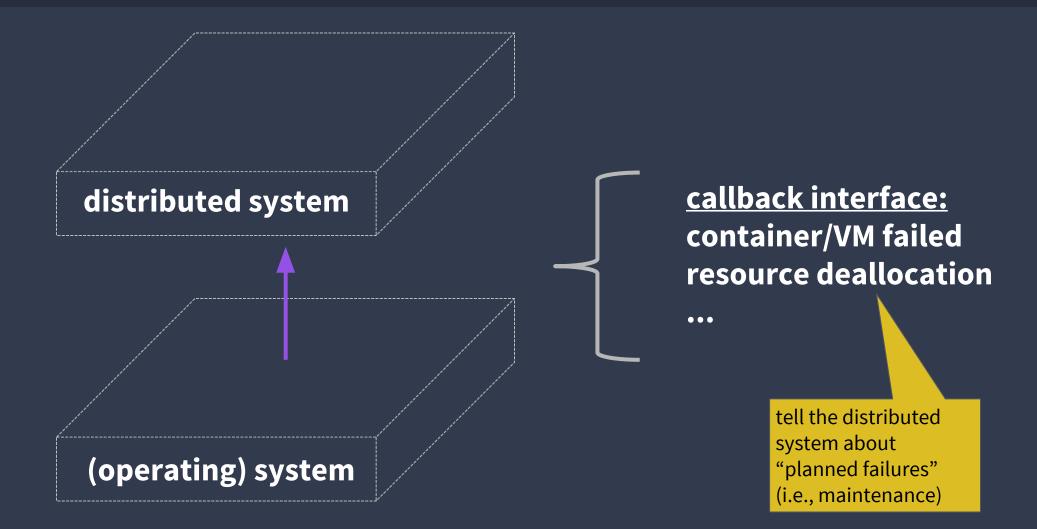
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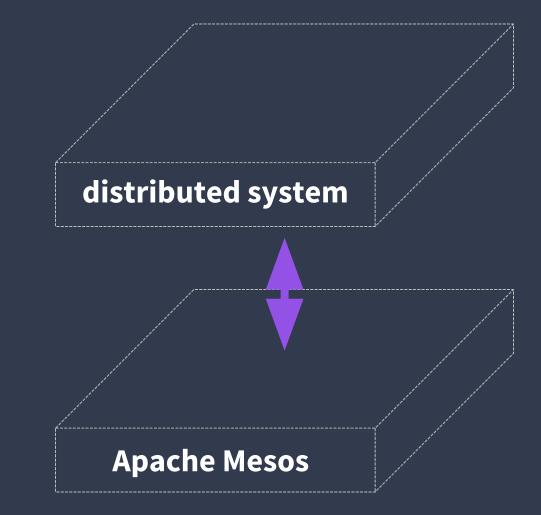
distributed systems need bidirectional interface too



distributed systems need bidirectional interface too



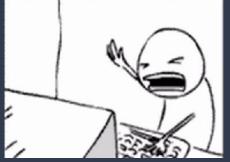
Apache Mesos



Dogfooding: Apache Spark



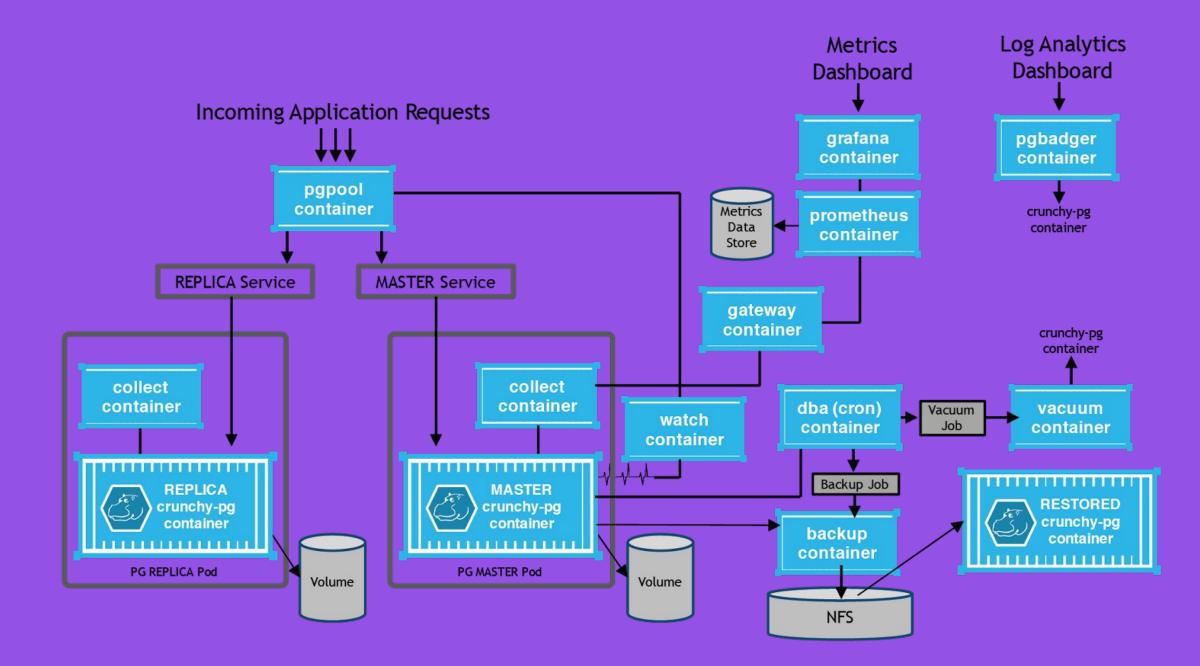
reality is people are (already) building software that operates distributed systems ...



common pattern: ad hoc control planes

goal: provide *distributed system** as software as a service (SaaS) to the rest of your internal organization or to sell to external organizations **solution:** a *control plane* built out of ad hoc scripts, ancillary services, etc, that deploy, maintain, and upgrade said SaaS

* e.g., analytics via Spark, message queue via Kafka, key/value store via Cassandra



\$ kubectl create -f \$LOC/kitchensink-master-service.json \$ kubectl create -f \$LOC/kitchensink-slave-service.json \$ kubectl create -f \$LOC/kitchensink-pgpool-service.json <u>\$ envsubst < \$LOC/kitchensink-sync-slave-pv.json | kubectl create -f -</u> \$ envsubst < \$LOC/kitchensink-master-pv.json | kubectl create -f -\$ kubectl create -f \$LOC/kitchensink-sync-slave-pvc.json \$ kubectl create -f \$LOC/kitchensink-master-pvc.json \$ envsubst < \$LOC/kitchensink-master-pod.json | kubectl create -f -\$ envsubst < \$LOC/kitchensink-slave-dc.json | kubectl create -f -\$ envsubst < \$LOC/kitchensink-sync-slave-pod.json | kubectl create -f -\$ envsubst < \$LOC/kitchensink-pgpool-rc.json | kubectl create -f -\$ kubectl create -f \$LOC/kitchensink-watch-sa.json \$ envsubst < \$LOC/kitchensink-watch-pod.json | kubectl create -f -

https://schd.ws/hosted_files/cnd2016/8d/Containerizing%20PostgreSQL%20and%20Making%20it%20Cloud%20Native%20Ready%20-%20Jeff%20McCormick.pdf

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what happens if there's a bug in the control plane?

what if my control plane has diverged from yours?

what happens when a new release of the distributed system invalidates an assumption the control plane previously made?

a better world ...

control planes should be built into the distributed systems itself by the experts who built the distributed system in the first place!

as an industry we should strive to build a standard interface that distributed systems can leverage

vice versa: abstractions exist for good reasons, but without sufficient communication they force sub-optimal outcomes ...

a better world ...

control planes should be built into distributed systems themselves by the experts who built the distributed system in the first place!

as an industry we should strive to build a standard interface distributed systems can leverage

our standard interface should be bidirectional to avoid sub-optimal outcomes

how do we scale the operations of distributed systems?

let them scale themselves!

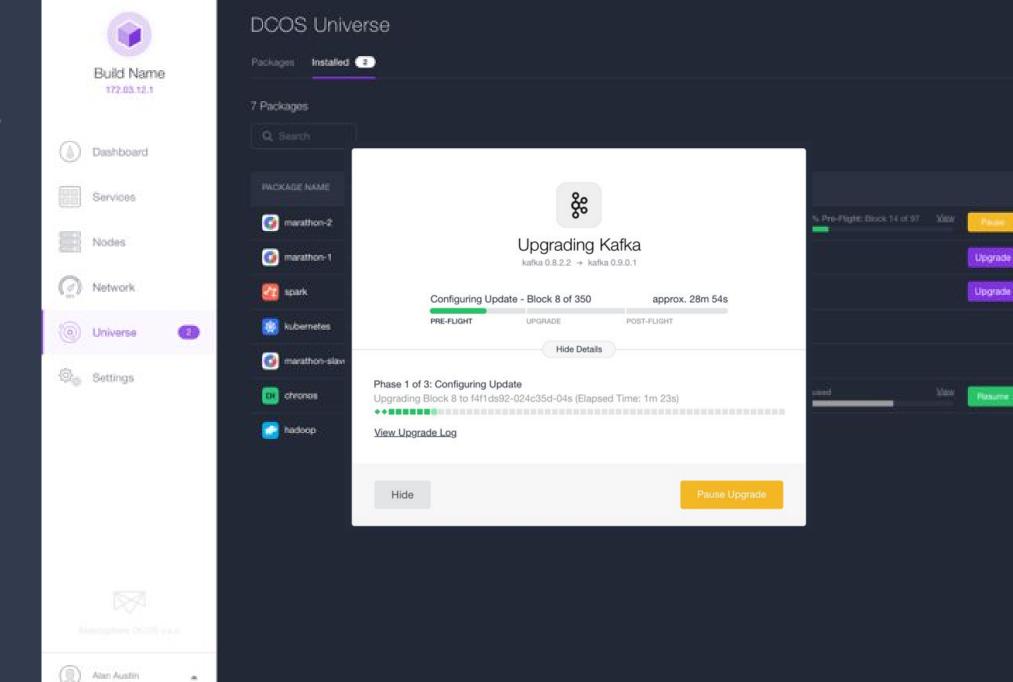
OPERATING SYSTEMS ARE FOR APPLICATIONS

"SaaS" Experience using DC/OS





DC/OS SERVICE MANAGES IT'S OWN UPGRADES



DC/OS: AVOIDING CLOUD LOCK-IN #2

	CAPABILITY	AWS	AZURE	GCP	DC/OS
Storage	Object Storage	S3	Blob Storage	Cloud Storage	Quobyte
	Block Storage	Elastic Block Storage (EBS)	Page Blobs, Premium Storage	GCE Persistent Disks	EMC ² ScaleiO
	File Storage	Elastic File System	File Storage	ZFS / Avere	EMC ² Scalei0
DB	Relational	RDS	SQL Database	Cloud SQL (MySQL)	MariaDB
	NoSQL	DynamoDB	DocumentDB	Datastore, Bigtable	cassandra ArangoDB Srick
Data & Analytics	Full Text Search	CloudSearch	Log Analytics, Search	N/A	😽 elastic
	Hadoop / Analytics	Elastic Map Reduce (EMR)	HDInsight	Dataproc, Dataflow	Spark
	Stream Processing / Ingest	Kinesis	Stream Analytics, Data Lake	Kinesis	kafka Spark
	Data Warehouse	Redshift	SQL Data Warehouse	BigQuery	
Other	Monitoring	CloudWatch	Application Insights, Portal	Stackdriver Monitoring	DATADOG
	Serverless	Lambda	Azure Functions	Google Cloud Functions	GALACTIC FOG

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THANK YOU!

DEMO!

QUESTIONS?

bigger picture: abstractions exist for good reasons, but without sufficient communication they force sub-optimal outcomes ...